

# A Java-like calculus with user-defined coeffects

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## Abstract

Modern applications are thought to be resource-aware, so it is very useful to focus on the concept of resource and to keep track of the use of them. *Coeffect systems* provide a static control capable to guarantee interesting properties on the usage of tracked objects. Our goal is to develop a Java-like calculus where declared variables can be annotated by *coeffects* specifying constraints on their use, such as linearity or security levels. Such annotations are written in the language itself, as expressions of type `Coeffect`, a predefined class which can be extended by user-defined subclasses, modeling coeffects desired for a specific application. We will also show a simple example of coeffect system checking the linear use of a variable.

## Coeffects

Coeffect systems are, in a sense, the dual of effect systems. The latter track *how the program modifies the environment*, coeffect systems *what the program requires from the context* of a computation. There are two kinds of coeffect systems:

- *Structural* coeffects annotate each variable in the context independently
- *Flat* coeffects annotates the whole context

Here we consider structural coeffects.

## Structure of coeffects

Coeffects are assumed to form a *semiring*, that is, a tuple  $(\mathcal{C}, +, 0, \times, 1)$  such that

- $(\mathcal{C}, +, 0)$  is a commutative monoid
- $(\mathcal{C}, \times, 1)$  is a monoid
- Given  $c$  in  $\mathcal{C}$
- Given  $c_1, c_2, c_3$  in  $\mathcal{C}$
- $-c_1 \times (c_2 + c_3) = (c_1 \times c_2) + (c_1 \times c_3)$
- $-(c_1 + c_2) \times c_3 = (c_1 \times c_3) + (c_2 \times c_3)$
- $-0 \times c = c \times 0 = 0$

## Coeffect annotations

Our Java-like calculus supports *user-defined* structural coeffects. That is, coeffect annotations are values of (subclasses of) a predefined class `Coeffect`, analogously to Java exceptions which are expressions of (subclasses of) `Exception`.

Coeffects can be seen as constraints on the use of declared variables, so with user-defined coeffects we can impose user-defined constraints. Metavariable  $\hat{v}$  is used for used-defined coeffects, that is, values expected to be of a subclass of `Coeffect`.

$$\{ T[\hat{v}] \ x = e; \ e' \}$$

This syntax imposes the constraint represented by  $\hat{v}$  on the use of  $x$  in  $e'$ .

Subclasses of `Coeffect` should implement methods:

- `sup`, corresponding to a  $\vee$  operator that induces an order relation among coeffects
- `sum` and `mult`, corresponding respectively to the  $+$  and  $\times$  operators of the semiring
- `zero` and `one`, returning a new `Coeffect` object corresponding respectively to 0 and 1

## Properties of user-defined coeffects

To guarantee preservation of coeffects during execution, operators determined by user-defined methods have to respect some equations, notably:

Given coeffects  $\hat{v}_1, \hat{v}_2, \hat{v}_3$ :

1.  $\hat{v}_1 + \hat{v}_2 = \hat{v}_2 + \hat{v}_1$
2.  $\hat{v}_1 + (\hat{v}_2 + \hat{v}_3) = (\hat{v}_1 + \hat{v}_2) + \hat{v}_3$
3.  $\hat{v}_1 + 0 = \hat{v}_1$
4.  $\hat{v}_1 \times 1 = 1 \times \hat{v}_1 = \hat{v}_1$
5.  $\hat{v}_1 \times 0 = 0 \times \hat{v}_1 = 0$
6.  $\hat{v}_1 \times (\hat{v}_2 + \hat{v}_3) = (\hat{v}_1 \times \hat{v}_2) + \hat{v}_1 \times \hat{v}_3$
7.  $\hat{v}_1 \times (\hat{v}_2 \times \hat{v}_3) = (\hat{v}_1 \times \hat{v}_2) \times \hat{v}_3$
8.  $(\hat{v}_1 + \hat{v}_2) \times \hat{v}_3 = \hat{v}_1 \times \hat{v}_3 + \hat{v}_2 \times \hat{v}_3$
9.  $\hat{v}_1 \vee \hat{v}_1 = \hat{v}_1$
10.  $\hat{v}_1 \vee (\hat{v}_2 \vee \hat{v}_3) = (\hat{v}_1 \vee \hat{v}_2) \vee \hat{v}_3$
11.  $\hat{v}_1 \vee \hat{v}_2 = \hat{v}_2 \vee \hat{v}_1$
12.  $\hat{v}_1 \vee (\hat{v}_1 + \hat{v}_2) = (\hat{v}_1 + \hat{v}_2)$
13.  $\hat{v}_1 + (\hat{v}_2 \vee \hat{v}_3) = (\hat{v}_1 + \hat{v}_2) \vee (\hat{v}_1 + \hat{v}_3)$
14.  $\hat{v}_1 \times (\hat{v}_2 \vee \hat{v}_3) = (\hat{v}_1 \times \hat{v}_2) \vee (\hat{v}_1 \times \hat{v}_3)$
15.  $(\hat{v}_2 \vee \hat{v}_3) \times \hat{v}_1 = (\hat{v}_2 \times \hat{v}_1) \vee (\hat{v}_3 \times \hat{v}_1)$

## An example: 0, 1, $\omega$ coeffects

In this simple example, coeffect `new Zero()` is meant to be assigned to unused variables, `new One()` to variables used linearly (exactly once), `new Omega()` to unrestricted variables. Class `Linearity` is used to define methods `zero` and `one` only once.

```
class Linearity extends Coeffect{
  Coeffect zero(){ new Zero()}
  Coeffect one(){new One()}
}

class Zero extends Linearity{
  Coeffect sup(Coeffect c) {
    case c of
      (Linearity x) x
      (Coeffect x) new Coeffect()
  }
  Coeffect sum(Coeffect c) {
    case c of
      (Linearity x) x
      (Coeffect x) new Coeffect()
  }
  Coeffect mult(Coeffect c) {
    case c of
      (Linearity x) new Zero()
      (Coeffect x) new Coeffect()
  }
}

class One extends Linearity{
  Coeffect sup(Coeffect c) {
    case c of
      (Zero x) new One()
      (One x) new One()
      (Omega x) new Omega()
      (Coeffect x)
        new Coeffect()
  }
  Coeffect sum(Coeffect c) {
    case c of
      (Zero x) new One()
      (One x) new Omega()
      (Omega x) new Omega()
      (Coeffect x)
        new Coeffect()
  }
  Coeffect mult(Coeffect c) {
    case c of
      (Linearity x) x
      (Coeffect x)
        new Coeffect()
  }
}

class Omega extends Linearity{
  Coeffect sup(Coeffect c) {
    case c of
      (Linearity x) new Omega()
      (Coeffect x) new Coeffect()
  }
  Coeffect sum(Coeffect c) {
    case c of
      (Linearity x) new Omega()
      (Coeffect x) new Coeffect()
  }
}

Coeffect mult(Coeffect c) {
  case c of
    (Zero x) new Zero()
    (One x) new Omega()
    (Omega x) new Omega()
    (Coeffect x)
      new Coeffect()
}
```

## Soundness result

Execution preserves types and coeffects assuming that user-defined coeffects guarantee conditions 1-15

## Future goals

- Adding *graded modal types*
- Allowing a “global” annotation in a method’s signature
- Allowing variables in coeffect annotations

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